



AIT 682: Network and Systems Security

Topic 5.1 Basic Number Theory -- Foundation of Public Key Cryptography

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Outline

- GCD and Euclid's Algorithm
- Modulo Arithmetic
- Modular Exponentiation
- Discrete Logarithms

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GCD and Euclid's Algorithm



Some Review: Divisors

- Set of all **integers** is $Z = \{\dots, -2, -1, 0, 1, 2, \dots\}$
- **b divides a** (or b is a *divisor* of a) if $a = mb$ for some m
 - denoted $b|a$
 - any $b \neq 0$ divides 0
- For any a , 1 and a are *trivial divisors* of a
 - all other divisors of a are called **factors** of a

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Primes and Factors

- a is **prime** if it has no non-trivial factors
 - examples: 2, 3, 5, 7, 11, 13, 17, 19, 31,...
- Theorem: there are infinitely many primes
- Any integer $a > 1$ can be factored in a unique way as $p_1^{a_1} \cdot p_2^{a_2} \cdot \dots \cdot p_t^{a_t}$
 - where all $p_1 > p_2 > \dots > p_t$ are prime numbers and where each $a_i > 0$

Examples:

$$91 = 13^1 \times 7^1$$

$$11011 = 13^1 \times 11^2 \times 7^1$$

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Common Divisors

- A number d that is a divisor of both a and b is a **common divisor** of a and b

Example: common divisors of 30 and 24 are 1, 2, 3, 6

- If $d|a$ and $d|b$, then $d|(a+b)$ and $d|(a-b)$

Example: Since $3|30$ and $3|24$, $3|(30+24)$ and $3|(30-24)$

- If $d|a$ and $d|b$, then $d|(ax+by)$ for any integers x and y

Example: $3|30$ and $3|24 \rightarrow 3|(2*30 + 6*24)$

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Greatest Common Divisor (GCD)

- $\text{gcd}(a,b) = \max\{k \mid k|a \text{ and } k|b\}$

Example: $\text{gcd}(60,24) = 12$, $\text{gcd}(a,0) = a$

- Observations
 - $\text{gcd}(a,b) = \text{gcd}(|a|, |b|)$
 - $\text{gcd}(a,b) \leq \min(|a|, |b|)$
 - if $0 \leq n$, then $\text{gcd}(an, bn) = n * \text{gcd}(a,b)$
- For all positive integers d , a , and b ...
 - ...if $d \mid ab$
 - ...and $\text{gcd}(a,d) = 1$
 - ...then $d \mid b$

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GCD (Cont'd)

- Computing GCD by hand:
 - if $a = p_1^{a_1} p_2^{a_2} \dots p_r^{a_r}$ and
 - $b = p_1^{b_1} p_2^{b_2} \dots p_r^{b_r}$,
 - ...where $p_1 < p_2 < \dots < p_r$ are prime,
 - ...and a_i and b_i are nonnegative,
 - ...then $\text{gcd}(a, b) =$
 - $p_1^{\min(a_1, b_1)} p_2^{\min(a_2, b_2)} \dots p_r^{\min(a_r, b_r)}$
- ⇒ Slow way to find the GCD
 - requires factoring a and b first (which can be slow)

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Euclid's Algorithm for GCD

- Insight:
 - $\text{gcd}(x, y) = \text{gcd}(y, x \bmod y)$
- Procedure **euclid(x, y)** :

```

r[0] = x, r[1] = y, n = 1;
while (r[n] != 0) {
    n = n+1;
    r[n] = r[n-2] % r[n-1];
}
return r[n-1];
    
```

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Example

n	r_n
0	595
1	408
2	$595 \bmod 408 = 187$
3	$408 \bmod 187 = 34$
4	$187 \bmod 34 = 17$
5	$34 \bmod 17 = 0$

$\gcd(595, 408) = 17$

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Running Time

- Running time is logarithmic in size of x and y
- Worst case occurs when ???

```

Enter x and y: 102334155 63245986
Step 1: r[i] = 39088169
Step 2: r[i] = 24157817
Step 3: r[i] = 14930352
Step 4: r[i] = 9227465
...
Step 34: r[i] = 5
Step 35: r[i] = 3
Step 36: r[i] = 2
Step 37: r[i] = 1
Step 38: r[i] = 0
gcd of 102334155 and 63245986 is 1
    
```

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Extended Euclid's Algorithm

- Let $\mathcal{LC}(x, y) = \{ux + vy : x, y \in \mathbb{Z}\}$ be the set of linear combinations of x and y
- Theorem: if x and y are any integers > 0 , then $\gcd(x, y)$ is the **smallest positive element of $\mathcal{LC}(x, y)$**
- Euclid's algorithm can be extended to **compute u and v** , as well as $\gcd(x, y)$
- Procedure **exteuclid**(x, y):
(next page...)

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Extended Euclid's Algorithm

```

r[0] = x, r[1] = y, n = 1;
u[0] = 1, u[1] = 0;
v[0] = 0, v[1] = 1;
while (r[n] != 0) {
    n = n+1;
    r[n] = r[n-2] % r[n-1];
    q[n] = (int) (r[n-2] / r[n-1]);
    u[n] = u[n-2] - q[n]*u[n-1];
    v[n] = v[n-2] - q[n]*v[n-1];
}
return r[n-1], u[n-1], v[n-1];

```

floor function

Exercise: Show $r[n]=u[n]x+v[n]y$

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Extended Euclid's Example

n	q_n	r_n	u_n	v_n
0	-	595	1	0
1	-	408	0	1
2	1	187	1	-1
3	2	34	-2	3
4	5	17	11	-16
5	2	0	-24	35

$$\gcd(595, 408) = 17 = 11 \cdot 595 + (-16) \cdot 408$$

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Extended Euclid's Example

n	q_n	r_n	u_n	v_n
0	-	99	1	0
1	-	78	0	1
2				
3				
4				
5				
6				

$$\gcd(99, 78) = 3 = -11 \cdot 99 + 14 \cdot 78$$

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Relatively Prime

- Integers a and b are *relatively prime* iff $\gcd(a,b) = 1$
 - example: 8 and 15 are relatively prime
- Integers n_1, n_2, \dots, n_k are *pairwise relatively prime* if $\gcd(n_i, n_j) = 1$ for all $i \neq j$

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Review of Modular Arithmetic

Remainders and Congruency

- For any integer a and any positive integer n , there are two unique integers q and r , such that $0 \leq r < n$ and $a = qn + r$
 - r is the *remainder* of division by n , written $r = a \bmod n$
- a and b are *congruent modulo n* , written $a \equiv b \pmod{n}$, if $a \bmod n = b \bmod n$

Example: $12 = 2 \cdot 5 + 2 \rightarrow 2 = 12 \bmod 5$

Example: $7 \bmod 5 = 12 \bmod 5 \rightarrow 7 \equiv 12 \pmod{5}$

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Negative Numbers

- In modular arithmetic,
...a negative number a is usually replaced
by the congruent number $b \pmod n$,
...where b is the smallest non-negative
number
...such that $b = a + m \cdot n$

Example: $-3 \equiv 4 \pmod 7$

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Remainders (Cont'd)

- For any positive integer n , the integers
can be divided into n equivalence classes
according to their remainders modulo n
– denote the set as Z_n
- i.e., the $(\pmod n)$ operator maps all
integers into the set of integers $Z_n = \{0, 1, 2, \dots, (n-1)\}$

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Modular Arithmetic

- Modular addition
– $[(a \pmod n) + (b \pmod n)] \pmod n = (a+b) \pmod n$
Example: $[16 \pmod{12} + 8 \pmod{12}] \pmod{12} = (16 + 8) \pmod{12} = 0$
- Modular subtraction
– $[(a \pmod n) - (b \pmod n)] \pmod n = (a - b) \pmod n$
Example: $[22 \pmod{12} - 8 \pmod{12}] \pmod{12} = (22 - 8) \pmod{12} = 2$
- Modular multiplication
– $[(a \pmod n) \times (b \pmod n)] \pmod n = (a \times b) \pmod n$
Example: $[22 \pmod{12} \times 8 \pmod{12}] \pmod{12} = (22 \times 8) \pmod{12} = 8$

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Properties of Modular Arithmetic

- **Commutative** laws
 - $(w + x) \bmod n = (x + w) \bmod n$
 - $(w \times x) \bmod n = (x \times w) \bmod n$
- **Associative** laws
 - $[(w + x) + y] \bmod n = [w + (x + y)] \bmod n$
 - $[(w \times x) \times y] \bmod n = [w \times (x \times y)] \bmod n$
- **Distributive** law
 - $[w \times (x + y)] \bmod n = [(w \times x) + (w \times y)] \bmod n$

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Properties (Cont'd)

- **Idempotent** elements
 - $(0 + m) \bmod n = m \bmod n$
 - $(1 \times m) \bmod n = m \bmod n$
 - **Additive** inverse $(-w)$
 - for each $m \in \mathbb{Z}_n$, there exists z such that $(m + z) \bmod n = 0$
 - alternatively, $z = (n - m) \bmod n$
- Example: 3 and 4 are additive inverses mod 7, since $(3 + 4) \bmod 7 = 0$
- **Multiplicative** inverse
 - for each positive $m \in \mathbb{Z}_n$ is there a z s.t. $m * z = 1 \bmod n$?

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Multiplicative Inverses

- Don't always exist!
 - Ex.: there is no z such that $6 \times z = 1 \bmod 8$

z	0	1	2	3	4	5	6	7	
$6 \times z$	0	6	12	18	24	30	36	42	...
$6 \times z \bmod 8$	0	6	4	2	0	6	4	2	

- An positive integer $m \in \mathbb{Z}_n$ has a multiplicative inverse $m^{-1} \bmod n$ iff $\gcd(m, n) = 1$, i.e., m and n are **relatively prime**
 - ⇒ If n is a prime number, then all positive elements in \mathbb{Z}_n have multiplicative inverses

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Finding the Multiplicative Inverse

- Given m and n , how do you find $m^{-1} \bmod n$?
- Extended Euclid's Algorithm
`exteuclid(m, n):`
 $m^{-1} \bmod n = v_{n-1}$
 - if $\gcd(m, n) \neq 1$ there is **no** multiplicative inverse $m^{-1} \bmod n$

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Example

n	q_n	r_n	u_n	v_n
0	-	35	1	0
1	-	12	0	1
2	2	11	1	-2
3	1	1	-1	3
4	11	0	12	-35

$$\gcd(35, 12) = 1 = -1 \cdot 35 + 3 \cdot 12$$

$$12^{-1} \bmod 35 = 3 \text{ (i.e., } 12 \cdot 3 \bmod 35 = 1)$$

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Modular Division

- If the inverse of $b \bmod n$ exists, then
 $(a \bmod n) / (b \bmod n) = (a * b^{-1}) \bmod n$

Example: $(13 \bmod 11) / (4 \bmod 11) = (13 * 4^{-1} \bmod 11) = (13 * 3) \bmod 11 = 6$

Example: $(8 \bmod 10) / (4 \bmod 10)$ not defined since 4 does not have a multiplicative inverse mod 10

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Modular Exponentiation (Power)

Modular Powers

Example: show the powers of $3 \pmod{7}$

i	0	1	2	3	4	5	6	7	8
3^i	1	3	9	27	81	243	729	2187	6561
$3^i \pmod{7}$	1	3	2	6	4	5	1	3	2

And the powers of $2 \pmod{7}$

i	0	1	2	3	4	5	6	7	8	9
2^i	1	2	4	8	16	32	64	128	256	512
$2^i \pmod{7}$	1	2	4	1	2	4	1	2	4	1

Fermat's "Little" Theorem

- If p is prime
 ...and a is a positive integer not divisible by p ,
 ...then $a^{p-1} \equiv 1 \pmod{p}$

Example: 11 is prime, 3 not divisible by 11,
 so $3^{11-1} = 59049 \equiv 1 \pmod{11}$

Example: 37 is prime, 51 not divisible by 37,
 so $51^{37-1} \equiv 1 \pmod{37}$

Useful?

Multiplicative Group Z_n^*

- Let Z_n^* be the set of numbers between 1 and $n-1$ that are **relatively prime to n**
- Z_n^* is **closed under multiplication mod n**
- Ex.: $Z_8^* = \{1, 3, 5, 7\}$

*	1	3	5	
1				
3				
5				
7				

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The Totient Function

- $\phi(n) = |Z_n^*|$ = the **number** of integers less than n and relatively prime to n
- a) if n is **prime**, then $\phi(n) = n-1$

Example: $\phi(7) = 6$

- b) if $n = p^\alpha$, where p is prime and $\alpha > 0$, then
 $\phi(n) = (p-1) \cdot p^{\alpha-1}$

Example: $\phi(25) = \phi(5^2) = 4 \cdot 5^1 = 20$

- c) if $n = p \cdot q$, and p, q are relatively prime, then
 $\phi(n) = \phi(p) \cdot \phi(q)$

Example: $\phi(15) = \phi(5 \cdot 3) = \phi(5) \cdot \phi(3) = 4 \cdot 2 = 8$

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Euler's Theorem

- For every a and n that are **relatively prime**,
 $a^{\phi(n)} \equiv 1 \pmod{n}$

Example: For $a = 3$, $n = 10$, which relatively prime:

$$\phi(10) = 4$$

$$3^{\phi(10)} = 3^4 = 81 \equiv 1 \pmod{10}$$

Example: For $a = 2$, $n = 11$, which are relatively prime:

$$\phi(11) = 10$$

$$2^{\phi(11)} = 2^{10} = 1024 \equiv 1 \pmod{11}$$

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More Euler..

- Variant:
for all n , $a^{k\phi(n)+1} \equiv a \pmod n$ for all a in \mathbb{Z}_n^* , and all non-negative k

Example: for $n = 20$, $a = 7$, $\phi(n) = 8$, and $k = 3$:

$$7^{3 \cdot 8 + 1} \equiv 7 \pmod{20}$$

- Generalized Euler's Theorem:
for $n = pq$ (p and q distinct primes),
 $a^{k\phi(n)+1} \equiv a \pmod n$ for all a in \mathbb{Z}_n^* and all non-negative k

Example: for $n = 15$, $a = 6$, $\phi(n) = 8$, and $k = 3$:

$$6^{3 \cdot 8 + 1} \equiv 6 \pmod{15}$$

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Modular Exponentiation

- $x^y \pmod n \equiv x^{y \pmod{\phi(n)}} \pmod n$

Example: $x = 5$, $y = 7$, $n = 6$, $\phi(6) = 2$

$$5^7 \pmod 6 = 5^{7 \pmod 2} \pmod 6 = 5 \pmod 6$$

- by this, if $y \equiv 1 \pmod{\phi(n)}$, then $x^y \pmod n \equiv x \pmod n$

Example:

$x = 2$, $y = 101$, $n = 33$, $\phi(33) = 20$, $101 \pmod{20} = 1$

$$2^{101} \pmod{33} = 2 \pmod{33}$$

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The Powers of An Integer, Modulo n

- Consider the expression $a^m \equiv 1 \pmod n$
- If a and n are relatively prime, then there is at least one integer m that satisfies the above equation
- Ex: for $a = 3$ and $n = 7$, what is m ?

i	1	2	3	4	5	6	7	8	9
$3^i \pmod 7$	3	2	6	4	5	1	3	2	6

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The Power (Cont'd)

- The **least** positive exponent m for which the above equation holds is referred to as...
 - the **order of $a \pmod n$** , or
 - the **length of the period generated by a**

Understanding Order of $a \pmod n$

- Powers of some integers a modulo 19

a	a^2	a^3	a^4	a^5	a^6	a^7	a^8	a^9	a^{10}	a^{11}	a^{12}	a^{13}	a^{14}	a^{15}	a^{16}	a^{17}	a^{18}	order	
1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1
2	4	8	16	13	7	14	9	18	17	15	11	3	6	12	5	10	1	18	
4	16	7	9	17	11	6	5	1	4	16	7	9	17	11	6	5	1	9	
7	11	1	7	11	1	7	11	1	7	11	1	7	11	1	7	11	1	3	
8	7	18	11	12	1	8	7	18	11	12	1	8	7	18	11	12	1	6	
9	5	7	6	16	11	4	17	1	9	5	7	6	16	11	4	17	1	9	
18	1	18	1	18	1	18	1	18	1	18	1	18	1	18	1	18	1	2	

Observations on The Previous Table

- The length of each period divides $18 = \phi(19)$
 - i.e., the lengths are 1, 2, 3, 6, 9, 18
- Some of the sequences are of length 18
 - e.g., the base **2** generates (via powers) all members of \mathbb{Z}_n^*
 - The base is called the **primitive root**
 - The base is also called the **generator** when n is prime

Reminder of Results

Totient function:

if n is prime, then $\phi(n) = n-1$

if $n = p^\alpha$, where p is prime and $\alpha > 0$, then $\phi(n) = (p-1) * p^{\alpha-1}$

if $n = p * q$, and p, q are relatively prime, then $\phi(n) = \phi(p) * \phi(q)$

Example: $\phi(7) = 6$

Example: $\phi(25) = \phi(5^2) = 4 * 5^1 = 20$

Example: $\phi(15) = \phi(5 * 3) = \phi(5) * \phi(3) = 4 * 2 = 8$

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Reminder (Cont'd)

- Fermat: If p is prime and a is positive integer not divisible by p , then $a^{p-1} \equiv 1 \pmod{p}$

Example: 11 is prime, 3 not divisible by 11, so $3^{11-1} = 59049 \equiv 1 \pmod{11}$

Euler: For every a and n that are relatively prime, then $a^{\phi(n)} \equiv 1 \pmod{n}$

Example: For $a = 3, n = 10$, which relatively prime: $\phi(10) = 4, 3^{\phi(10)} = 3^4 = 81 \equiv 1 \pmod{10}$

Variant: for all a in \mathbb{Z}_n^* , and all non-negative $k, a^{k\phi(n)+1} \equiv a \pmod{n}$

Example: for $n = 20, a = 7, \phi(n) = 8$, and $k = 3: 7^{3*8+1} \equiv 7 \pmod{20}$

Generalized Euler's Theorem: for $n = pq$ (p and q are distinct primes), all a in \mathbb{Z}_n^* , and all non-negative $k, a^{k\phi(n)+1} \equiv a \pmod{n}$

Example: for $n = 15, a = 6, \phi(n) = 8$, and $k = 3: 6^{3*8+1} \equiv 6 \pmod{15}$

$x^y \pmod{n} = x^{y \pmod{\phi(n)}} \pmod{n}$

Example: $x = 5, y = 7, n = 6, \phi(6) = 2, 5^7 \pmod{6} = 5^{7 \pmod{2}} \pmod{6} = 5^1 \pmod{6} = 5 \pmod{6}$

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Computing Modular Powers Efficiently

- The repeated squaring algorithm for computing $a^b \pmod{n}$
- Let b_i represent the i^{th} bit of b (total of k bits)

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Computing (Cont'd)

Algorithm `modexp(a, b, n)`

```

d = 1;
for i = k downto 1 do
  d = (d * d) % n; /* square */
  if (b_i == 1)
    d = (d * a) % n; /* step 2 */
  endif
enddo
return d;
  
```

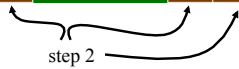
at each iteration, not just at end
 Requires time $\propto k = \text{logarithmic in } b$

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Example

- Compute $a^b \pmod n = 7^{560} \pmod{561} = 1 \pmod{561}$

i	10	9	8	7	6	5	4	3	2	1
b_i	1	0	0	0	1	1	0	0	0	0
d	1	7	49	157	526	160	241	298	166	67



Q: Can some other result be used to compute this particular example more easily? (Note: $561 = 3 \cdot 11 \cdot 17$.)

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Discrete Logarithms

Square Roots

- x is a *non-trivial square root of 1 mod n* if it satisfies the equation $x^2 \equiv 1 \pmod{n}$, but x is neither 1 nor $-1 \pmod{n}$

Ex: 6 is a square root of 1 mod 35 since $6^2 \equiv 1 \pmod{35}$

- Theorem: if there exists a non-trivial square root of 1 mod n , then n is **not** a prime
 - i.e., prime numbers will not have non-trivial square roots

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Roots (Cont'd)

- If $n = 2^{\alpha_0} p_1^{\alpha_1} p_2^{\alpha_2} \dots p_k^{\alpha_k}$, where $p_1 \dots p_k$ are distinct primes > 2 , then the **number of square roots** (including trivial square roots) are:
 - 2^k if $\alpha_0 \leq 1$

Example: for $n = 70 = 2^1 * 5^1 * 7^1$, $\alpha_0 = 1$, $k = 2$, and the number of square roots $= 2^2 = 4$ (1,29,41,69)

– 2^{k+1} if $\alpha_0 = 2$

Example: for $n = 60 = 2^2 * 3^1 * 5^1$, $k = 2$, the number of square roots $= 2^3 = 8$ (1,11,19,29,31,41,49,59)

– 2^{k+2} if $\alpha_0 > 2$

Example: for $n = 24 = 2^3 * 3^1$, $k = 1$, the number of square roots $= 2^3 = 8$ (1,5,7,11,13,17,19,23)

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Primitive Roots

- Reminder: the highest possible order of $a \pmod{n}$ is $\phi(n)$
- If the **order of $a \pmod{n}$** is $\phi(n)$, then a is referred to as a *primitive root of n*
 - for a prime number p , if a is a primitive root of p , then a, a^2, \dots, a^{p-1} are all distinct numbers mod p
- No simple general formula to compute primitive roots modulo n
 - there are methods to locate a primitive root faster than trying out all candidates

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Primitive Roots (Cont'd)

- Theorem: the **only** integers with primitive roots are of the form **2**, **4**, p^α , and $2p^\alpha$, where
 - p is any prime > 2
 - α is a positive integer

Example: for $n = 4$, $\phi(n) = 2$, primitive roots = $\{3\}$

Example: for $n = 3^2 = 9$, $\phi(n) = 6$, primitive roots = $\{2, 5\}$

Example: for $n = 19$, $\phi(n) = 18$, primitive roots = $\{2, 3, 10, 13, 14, 15\}$

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Discrete Logarithms

- For a primitive root a of a number p , where
 - $a^i \equiv b \pmod{p}$, for some $0 \leq i \leq p-1$
 - the exponent i is referred to as *the index of b for the base $a \pmod{p}$* , denoted as $\text{ind}_{a,p}(b)$
 - i is also referred to as the *discrete logarithm of b to the base $a, \text{ mod } p$*

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Logarithms (Cont'd)

- Example: 2 is a primitive root of 19.
The powers of 2 mod 19 =

b	1	2	3	4	5	6	7	8	9
$\text{ind}_{2,19}(b) = \log(b) \text{ base } 2 \text{ mod } 19$	0	1	13	2	16	14	6	3	8

10	11	12	13	14	15	16	17	18
17	12	15	5	7	11	4	10	9

Given a , i , and p , computing $b = a^i \pmod{p}$ is straightforward

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Computing Discrete Logarithms

- However, given a , b , and p , computing $i = \text{ind}_{a,p}(b)$ is **difficult**
 - Used as the basis of some public key cryptosystems

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Computing (Cont'd)

- Some properties of discrete logarithms
 - $\text{ind}_{a,p}(1) = 0$ because $a^0 \bmod p = 1$ *warning: $\phi(p)$, not p !*
 - $\text{ind}_{a,p}(a) = 1$ because $a^1 \bmod p = a$
 - $\text{ind}_{a,p}(yz) = (\text{ind}_{a,p}(y) + \text{ind}_{a,p}(z)) \bmod \phi(p)$

Example: $\text{ind}_{2,19}(5*3) = (\text{ind}_{2,19}(5) + \text{ind}_{2,19}(3)) = 11 \bmod 18$

– $\text{ind}_{a,p}(y^r) = (r \text{ind}_{a,p}(y)) \bmod \phi(p)$

Example: $\text{ind}_{2,19}(3^3) = (3 * \text{ind}_{2,19}(3)) = 3 \bmod 18$

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More on Discrete Logarithms

- Consider:
 - $x = a^{\text{ind}_{a,p}(x)} \bmod p$, Ex: $3 = 2^{13} \bmod 19$
 - $y = a^{\text{ind}_{a,p}(y)} \bmod p$, and Ex: $5 = 2^{16} \bmod 19$
 - $xy = a^{\text{ind}_{a,p}(xy)} \bmod p$ Ex: $3*5 = 2^{11} \bmod 19$
- 1) $a^{\text{ind}_{a,p}(xy)} \bmod p = (a^{\text{ind}_{a,p}(x)} \bmod p)(a^{\text{ind}_{a,p}(y)} \bmod p)$
Ex: $15 = 3 * 5$
 - 2) $a^{\text{ind}_{a,p}(xy)} \bmod p = (a^{\text{ind}_{a,p}(x) + \text{ind}_{a,p}(y)}) \bmod p$
Ex: $15 = 2^{13+16} \bmod 19$
 - 3) by Euler's theorem: $a^z = a^q \bmod p$ iff $z = q \bmod \phi(p)$

Ex: $15 = 2^{11} \bmod 19 = 2^{29} \bmod 19 \Leftrightarrow 11 = 29 \bmod 18$

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Summary

1. Number theory is the basis of public key cryptography
2. Euclid's algorithm is used to find GCD and multiplicative inverse
3. Computing $a^b \pmod n$ is accomplished by repeated squaring
4. Only primes have discrete logarithms, and they are expensive to compute
